# A Theory of Stable Market Segmentations

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August 29, 2024

- Homeowners associations
- Student groups
- Employer-based prices

Consumers interact with a seller as a group

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- Leads to "market segmentation"
- A cooperative approach:
  - segment = a coalition of consumers
- "Stable" segmentations
  - have good welfare properties













 $(C_1, 1)$ : a segment  $(C_1, 2)$ : a segment  $(C_2, 1)$ : not a segment  $(C_2, 2)$ : a segment



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$$CS(c,S) = \begin{cases} \max\{v(c) - 1, 0\} & \text{if } c \in C_1, \\ \max\{v(c) - 2, 0\} & \text{if } c \in C_2. \end{cases}$$



## Outline



### Stability

Definition (Objection)

A segment (C, p) objects to segmentation S if

 $\max\{v(c) - p, 0\} > CS(c, S)$  for some (measure > 0)  $c \in C$ 

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**Note:** Objecting segment  $(C, p) \notin S$ 

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A segment (C, p) objects to segmentation S if

 $\begin{array}{ll} \max\{v(c)-p,0\}\geq CS(c,S) \text{ for all } & c\in C\\ \max\{v(c)-p,0\}>CS(c,S) \text{ for some (measure }>0) \ c\in C \end{array}$ 

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Definition (Core)

S is in the core if  $\nexists$  segment (C, p) that objects to S

Two type illustration: Core may be empty



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# Two type illustration: Core may be empty If $\delta < 0.8$ : $S = \{(C_1, 1), (C_2, 2)\}$ not in core If $\delta = 0.8$ : $S = \{(C_1, 1), (C_2, 2)\}$ not in core

Segment  $(C'_1, 1)$  objects





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S is in weakened core if  $\nexists$  segment (C', p') that strictly objects to S:



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#### Definition (Weakened Core)

S is in weakened core if  $\nexists$  segment (C', p') that strictly objects to S:

- (C', p') objects to S
- ∀(C, p) ∈ S that is "broken apart" by C' (C ∩ C' and C\C' have positive measures), ∃ consumers in C ∩ C' who strictly benefit (CS(c, p') > CS(c, S))



### Cost-based justification



## Cost-based justification

- Small fee to break up existing coalitions
- Must be paid by members who want to deviate (in  $C_1 \cap C'_1$ )
- If not paid, objection fails



Two type illustration and weakened core  $S = \{(C_1, 1), (C_2, 2)\}$  is in the weakened core



Two type illustration and weakened core  $S = \{(C_1, 1), (C_2, 2)\}$  is in the weakened core  $S' = \{(C'_1, 2)\}$  is not in the weakened core



# Characterization

### Proposition

For any segmentation S, the following are equivalent

- **1** *S* is in the weakened core
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# Characterization

### Proposition

For any segmentation S, the following are equivalent

- **1** *S* is in the weakened core
- **2** the induced canonical segmentation of S efficient and saturated
- **3** S is "stable":  $\forall S' \neq S$ ,  $\exists (C, p) \in S$  that objects to S'

 $S = \{(C_1, 1), (C_2, 2)\}$ 















Is defined recursively. Let  $\bar{C} = [0,1]$ ,  $S = \emptyset$ 

- C := largest coalition where all prices (among remaining values in C
  are revenue-maximizing
- 2 Add  $(C, \underline{v}(C))$  to S
- Remove C from  $\bar{C}$
- Repeat until  $\bar{C} = \emptyset$

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The MER segmentation is stable

Bergemann, Brooks, Morris (2015):

- The MER segmentation maximizes consumer surplus
- But is not the only one









Thanks!

Stability ⇒ maximizing consumer surplus

Stability  $\Rightarrow$  maximizing consumer surplus



### Stability $\Rightarrow$ maximizing consumer surplus

 $S = \{(C_1, 1), (C_2, 3)\}$  is efficient and saturated  $\Rightarrow$  stable



Stability  $\notin$  maximizing consumer surplus

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### Stability $\not\leftarrow$ maximizing consumer surplus

- $S = \{(C_1, 1), (C_2, 2)\}$  maximizes consumer surplus
  - Efficient allocation
  - price 3 is revenue-maximizing for  $C_1, C_2, [0, 1]$



### Stability $\not\leftarrow$ maximizing consumer surplus

 $S = \{(C_1, 1), (C_2, 2)\}$  maximizes consumer surplus

- Efficient allocation
- ▶ price 3 is revenue-maximizing for  $C_1, C_2, [0, 1]$

S is not saturated and so not stable:

