Black-Box Credit Scoring and Data Sharing

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Introduction

Predictive algorithms are used in economic transactions to solve asymmetric information problems.

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Loan Approvals:

- FinTechs 3-1-0 model (Ant Group, Prosper, Lending Club, Kaggle...),
- Traditional lenders (JP Morgan, Bank of America...),
- Credit bureaus (UltraFico, FICO X Data, Vantage Score...).

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Loan Approvals:

Input Data \longrightarrow Algorithm \longrightarrow Allocation z $a_{\lambda}(z)$ $\ell \in \{0, 1\}$ $x \in \mathbb{R}$

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Loan Approvals:



Transparency vs Opacity:

- Opacity: users hide from the system (privacy-preserving apps, not sharing cookies, paying cash, ...),
- Transparency: users game the system (YouTube tutorials to increase FICO score, ...).

This Paper

Question: Should credit scoring algorithms be transparent or opaque? (EU Artificial Intelligence Act)

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- Opacity induces *hedging* : withholding is a safe strategy against unpredictability,
 - all evidence is withheld when the lender has bargaining power,
 - most conclusive evidence is disclosed when the borrower has bargaining power,

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- Opacity induces *hedging* : withholding is a safe strategy against unpredictability,
 - all evidence is withheld when the lender has bargaining power,
 - most conclusive evidence is disclosed when the borrower has bargaining power,
- The lender's optimal transparency regime:
 - is driven by credit rationing motives,
 - depends on the lender's bargaining power,
 - is often socially inefficient.



\mathbf{Model}







• $\theta \sim U[0,1]$ is unknown (s.t. $\mathbb{E}(\theta) > 1/X$),

 $\mu_{\lambda}(z) \triangleq \mathbb{E}(\theta|z) = \lambda \, z + (1-\lambda) \, \frac{1}{2},$

• can be estimated from data $z \in [0, 1]$:

where
$$\lambda \in [-1, 1]$$
.

DGP



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Allocation Algorithm: An algorithm is $a_{\lambda}(z) = (\ell_{\lambda}(z), x_{\lambda}(z))$, where

- $\ell \in \{0, 1\}$ credit provision decision,
- $x \in \mathbb{R}$ gross interest payment.



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Two-Sided Private Information:

- z is the borrower's private information,
- λ is the lender's private information.



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Optimal Credit Scoring

Lemma 1 (Optimal Credit Scoring)

If the lender knows the borrower's data point z, the optimal allocation algorithm is:

$$x_{\lambda}(z) = \frac{1}{\mu_{\lambda}(z)} + \phi\left(X - \frac{1}{\mu_{\lambda}(z)}\right),$$
$$\ell_{\lambda}(z) = \mathbb{1}\left\{\mu_{\lambda}(z) \ge \frac{1}{X}\right\}.$$





















Data-Sharing to a Transparent Algorithm (Gaming)



















Proposition 1 (Data-Sharing to a Transparent Algorithm)

With a transparent algorithm, the set of borrowers withholding data is:

$$\mathcal{G}(\lambda,\pi) = \begin{cases} \begin{bmatrix} 0, \max\left\{r(\lambda); \gamma(\pi)\right\} \end{bmatrix} & \text{if } \lambda \in (0,1], \\ \begin{bmatrix} 0,1 \end{bmatrix} & \text{if } \lambda = 0, \\ \begin{bmatrix} \min\left\{r(\lambda); 1 - \gamma(\pi)\right\}, 1 \end{bmatrix} & \text{if } \lambda \in [-1,0) \end{cases}$$

where $\gamma(\pi)$ is increasing.

Data-Sharing to an Opaque Algorithm (Hedging)























Opacity \Rightarrow Hedging (Case 2: Competitive Markets)



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Proposition 2 (Data-Sharing to an Opaque Algorithm)

With an opaque algorithm, the set of borrowers withholding data is:

$$\mathcal{H}(b,\phi) = \begin{cases} [0,1] & \text{if } \phi \ge 1-2b, \\ [\eta(b,\phi), 1-\eta(b,\phi)] & \text{if } \phi < 1-2b, \end{cases}$$

where $\eta(b, \phi)$ is decreasing in ϕ and b.

Transparency vs Opacity

Transparency vs Opacity: Misallocations



Transparency vs Opacity: Misallocations



Transparency vs Opacity: Lender's Profits

Proposition 3 (Transparency vs Opacity: Lender's Profits)

The lender chooses an opaque algorithm iff $\pi > \hat{\pi}(b, \phi)$.



 ϕ (Lender's Bargaining Power)

Conclusion

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In the Paper:

- Redistributive effects,
- Welfare consequences.



FinTech.

Agarwal et al.(2019), Berg et al.(2020), Gambacorta et al.(2020), Fuster et al.(2022). **This Paper:** Opaque statistical technologies have welfare impact.

Voluntary Disclosure / Economics of Data and Privacy.

Grossman (1981), Milgrom (1981), Dye (1985), Quigley and Walther (2022), Bond and Zeng (2022); Ali, Lewis and Vasserman (2023), He, Huang and Zhou (2023).

This Paper: No unravelling when algorithm is opaque, Comparison.

Strategic Classification.

Frankel and Kartik (2019), Frankel and Kartik (2021), Ball (2022), Perez-Richet and Skreta (2022). This Paper: Opacity can soften gaming, but induce hedging.

Economics and Regulation of Algorithms.

Kleinberg et al. (2018), Cowgill and Tucker (2020), Cowgill and Stevenson (2020), Liang et al. (2023). This Paper: Optimal transparency regime.

Data-Generating Process

The DGP is truth-or-noise (Lewis and Sappington (1994)) allowing for negative correlation:

$$z = \begin{cases} \theta & \text{with pr. } \lambda \\ \varepsilon & \text{with pr. } 1 - \lambda \end{cases} \quad \text{if } \lambda \ge 0,$$

$$z = \begin{cases} 1 - \theta & \text{with pr.} \quad |\lambda| \\ \varepsilon & \text{with pr.} \quad 1 - |\lambda| \end{cases} \quad \text{if } \lambda < 0,$$

where $\varepsilon \sim U[0,1]$ is noise independent of θ .

The conditional pdf is:

$$f_{\lambda}(\theta|z) = \begin{cases} \lambda \,\delta(\theta-z) + (1-\lambda)1/2 & \text{if } \lambda \ge 0\\ \lambda \,\delta(1-\theta-z) + (1-\lambda)1/2 & \text{if } \lambda < 0, \end{cases}$$

where $\delta(\theta - z)$ is the Dirac's delta function.



Payoffs



$$V^{B}(\theta) = \ell ((\theta(X - x) + b))$$
$$V^{L}(\theta) = \ell (\theta x - 1)$$
$$W(\theta) = \ell (\theta X - 1 + b)$$

Back

Lemma A.1 (Credit Allocation following Data Withholding)

Let $Q \triangleq \{z \in [0,1] : m(z) = \emptyset\}$ be the set of borrowers that withhold data, the optimal credit allocation is:

$$x_{\lambda}(\varnothing) = \frac{1}{\mu_{\lambda}(\varnothing)} + \phi\left(X - \frac{1}{\mu_{\lambda}(\varnothing)}\right),$$

$$\ell_{\lambda}(\varnothing) = \mathbb{1}\left\{\mu_{\lambda}(\varnothing) > \frac{1}{X}\right\}.$$

where

$$\mu_{\lambda}(\emptyset) = \lambda \, z(\pi, Q) + (1 - \lambda) \frac{1}{2},$$

$$z(\pi, Q) = \omega(\pi, Q)\frac{1}{2} + (1 - \omega(\pi, Q))\mathbb{E}(\theta|\theta \in Q),$$
$$\omega(\pi, Q) = \frac{\pi}{\pi + (1 - \pi)\operatorname{Pr}(z \in Q)}.$$

