

Competitive Search and the Social Value of Public Information

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Motivation

- **Policy uncertainty/communication** is important
 - monetary surprises (Lucas (1972, 1973, 1975))
 - global games (Moris and Shin (2002))
 - usually: **frictionless trade**
- **Competitive search** has strong efficiency properties
 - market prices the good and the likelihood of trade
 - usually: market conditions are known/**no aggregate risk**

This paper

- Competitive (posting + commitment + directed) search,
- aggregate risk, and
- public information.

- Characterize equilibria.
 - What does aggregate risk imply for competitive search equilibrium?
 - How do search frictions interact with information friction?
- Study the effects of information on welfare.

Results

- **Arbitrary** marginal effects of public information.
- **Price dispersion** even with identical sellers and buyers.
- Market freezes (no trade in some states).
- Less efficient single-price than price-dispersed equilibria.
- Entry is generally inefficient.
- When/if more general mechanism improves upon price posting.

Related Literature

- Price setting under incomplete information (e.g. Keller and Rady, 1999, Hellwig and Venkateswaran, 2009): [frictionless trading](#)
- Search and aggregate uncertainty & incomplete information: Muring (2017), Lauermaun et al. (2018), Shneyerov and Wong (2020): [random search](#)
- Competitive search (Moen, 1997) with incomplete information (e.g. Guerrieri et al., 2010, Moen and Rosén, 2011, Julien and Roger, 2019, and Mayr-Dorn, 2020): [uncertainty about private/individual state](#)

Model

Sellers and buyers

- Competitive search model with **unknown demand state**.
- Sellers
 - fixed mass $\mathcal{S} = 1$,
 - **one unit** of indivisible good each,
 - post prices (can mix).
- Buyers
 - mass \mathcal{B}_i ,
 - unit demand,
 - each values good at v_i ,
 - **see prices** and decide which firm to contact (can mix),
 - trade off **price** and **probability of getting good**.

Matching

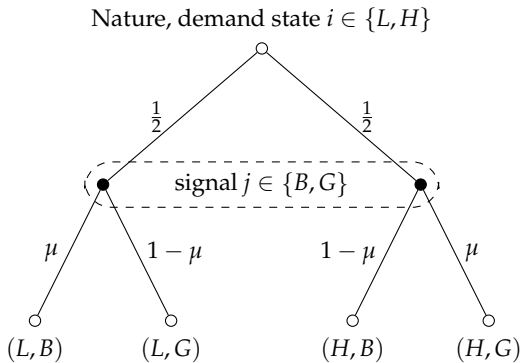
- Matching function $M(\mathcal{B}, \mathcal{S})$.
- Buyer-seller ratio: $x = \frac{\mathcal{B}}{\mathcal{S}}$.
- Probability of selling: $\lambda(x)$ with $\lambda' > 0, \lambda'' < 0$.
- Probability of buying: $\eta(x)$ with $\eta' < 0, \eta'' > 0$.
- Bilateral meetings.

Uncertainty and information

- **Uncertainty about state of demand** $i \in \{L, H\}$ with (\mathcal{B}_i, v_i)
 - uncertainty about buyer-seller ratio: $\mathcal{B}_H \geq \mathcal{B}_L$, (tightness uncertainty)
 - uncertainty about valuation: $v_H \geq v_L$ (surplus uncertainty)
 - Today: $\mathcal{B}_H > \mathcal{B}_L$ and $v_H = v_L = v$.
- **Information**
 - Buyers know state.
 - Sellers get **public signal** $j \in \{G, B\}$ before setting prices

$$\mu = Pr(j = G|i = H) = P(j = B|i = L) \in \left[\frac{1}{2}, 1 \right].$$

States and signals



Timing

- 1 Nature draws state $i \in \{H, L\}$.
- 2 Public signal is realised and sellers see outcome $j \in \{G, B\}$.
- 3 Sellers post prices.
- 4 Buyers contact sellers.
- 5 A buyer can buy at each seller that meets at least one buyer.
- 6 Trade.
- 7 Utilities are realised.

Equilibrium

- Strategies
 - firms see signal and post prices (can mix),
 - buyers see prices and choose prices to contact (can mix).
- Equilibrium
 - optimal prices p^j for firms, $j = \{G, B\}$,
 - contacting probabilities for buyers: buyer-seller ratios x_i^j , $i = \{H, L\}$ and $j = \{G, B\}$,
 - market clearing: buyer-seller ratios consistent with total measures of buyers and sellers.
- Symmetric equilibria.

Perfect information

Perfect information problem

- Sellers **know state** is $i \in \{L, H\}$.
- Sellers compete against market utility of buyers $U_i \leq v$ in state i :

$$\max_{p_i} \pi_i(p_i) := \lambda(x_i)p_i,$$

$$\text{subject to: } \eta(x_i)(v - p_i) - U_i = 0.$$

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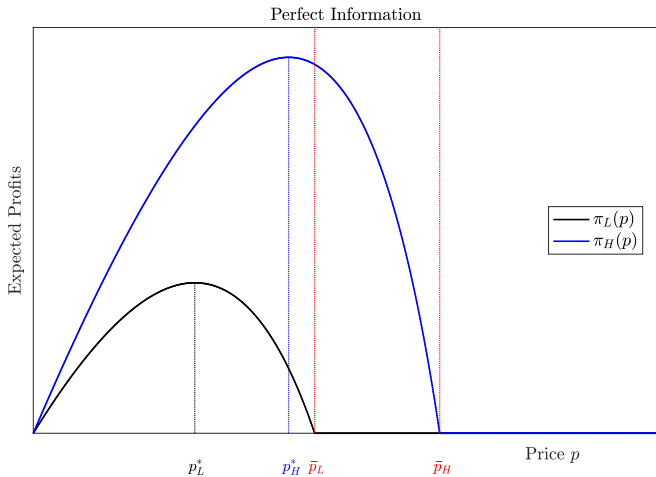
$$\begin{aligned} \max_{p_i} \pi_i(p_i) &:= \lambda(x_i)p_i, \\ \text{subject to: } &\eta(x_i)(v - p_i) - U_i = 0. \end{aligned}$$

- Solution:

$$p_i^* = \phi(x_i)v.$$

- $\phi(x_i) = -\frac{x_i \eta'(x_i)}{\eta(x_i)}$: elasticity of buying probability.
- Perfect information: **single price** in each state.

Perfect information profits



Imperfect information

Imperfect information problem

- Sellers **do not know state** $i \in \{L, H\}$, see $j = G$ or B .
- Seller who sees $j = G$ chooses p^G to

$$\max_{p^G} \pi^G(p^G) := \left[\mu \lambda \left(x_H^G \right) + (1 - \mu) \lambda \left(x_L^G \right) \right] p^G \quad (1)$$

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- If p^G is acceptable in state i , $x_i^G > 0$.
- If p^G is unacceptable to buyers in state i , $x_i^G = 0$.

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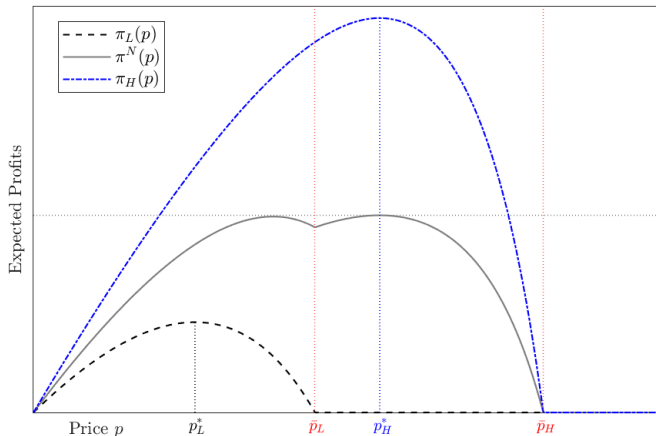
- If p^G is acceptable in state i , $x_i^G > 0$.
- If p^G is unacceptable to buyers in state i , $x_i^G = 0$.
- Similar problem for sellers who see $j = B$.

Pricing under imperfect information

- Pricing for **both states**, or pricing for **high state only**. [▶ Details](#)

Pricing under imperfect information

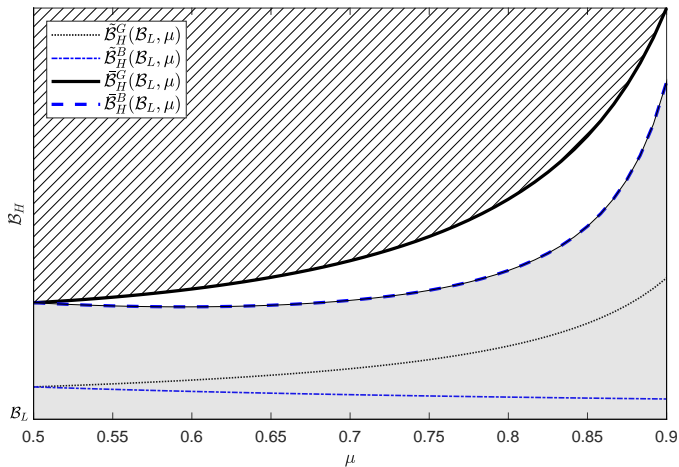
- Pricing for both states, or pricing for high state only. [▶ Details](#)



Theorem (Equilibria, tightness risk)

- ① There exist thresholds $\tilde{\mathcal{B}}_H^j := \tilde{\mathcal{B}}_H^j(\mathcal{B}_L, \mu)$, $\mathcal{B}_L < \tilde{\mathcal{B}}_H^j < \infty$ such that for $\mathcal{B}_H \in (\mathcal{B}_L, \tilde{\mathcal{B}}_H^j]$ the equilibrium profit function $\pi^j(p)$ is unimodal. Furthermore, $\partial \tilde{\mathcal{B}}_H^j / \partial \mathcal{B}_L > 0$, $\partial \tilde{\mathcal{B}}_H^G / \partial \mu > 0$, $\partial \tilde{\mathcal{B}}_H^B / \partial \mu < 0$ and $\tilde{\mathcal{B}}_H^G \geq \tilde{\mathcal{B}}_H^B$ with equality only if $\mu = 1/2$,
- ② there exist thresholds $\bar{\mathcal{B}}_H^j := \bar{\mathcal{B}}_H^j(\mathcal{B}_L, \mu)$; $\tilde{\mathcal{B}}_H^j < \bar{\mathcal{B}}_H^j \leq \infty$ such that for $\mathcal{B}_H \in (\tilde{\mathcal{B}}_H^j, \bar{\mathcal{B}}_H^j)$ the equilibrium profit function $\pi^j(p)$ is bimodal, but pricing for both states maximizes profits. Furthermore, $\lim_{\mu \rightarrow 1} \tilde{\mathcal{B}}_H^j(\mathcal{B}_L, \mu) = \infty$, and $\bar{\mathcal{B}}_H^j(\mathcal{B}_L, \mu) < \infty$ for \mathcal{B}_L and μ small enough,
- ③ a unique PSE exists iff $\mathcal{B}_H \leq \bar{\mathcal{B}}_H^j(\mathcal{B}_L, \mu)$ and a unique MSE exists iff $\mathcal{B}_H > \tilde{\mathcal{B}}_H^j(\mathcal{B}_L, \mu)$.

Equilibria



Both signals mix in hatched and post single price in shaded area.

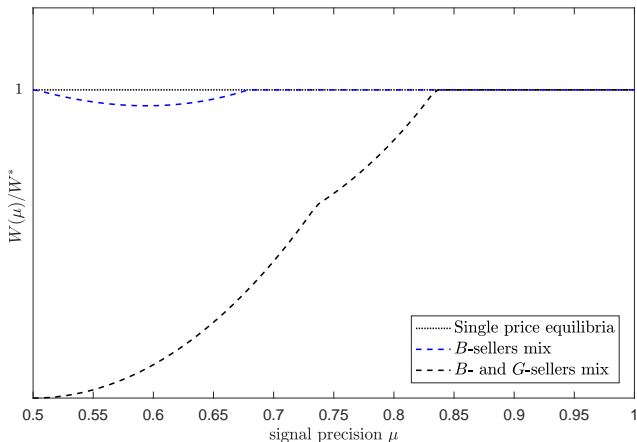
Welfare measure

- We measure welfare as expected value of trades:

$$W(\mu) = \frac{1}{2} \left[\mu \sum_{k=1}^{K^G} \kappa^{G,k} \lambda(x_H^{G,k}) + (1 - \mu) \sum_{k=1}^{K^B} \kappa^{B,k} \lambda(x_H^{B,k}) \right] v_H \\ + \frac{1}{2} \left[\mu \sum_{k=1}^{K^B} \kappa^{B,k} \lambda(x_L^{B,k}) + (1 - \mu) \sum_{k=1}^{K^G} \kappa^{G,k} \lambda(x_L^{G,k}) \right] v_L$$

- Price level does not matter.
- Price **dispersion** is **inefficient**.

Information and welfare



Normalised welfare for three different \mathcal{B}_H (fixed v and \mathcal{B}_L).

Conclusions

- Competitive search and **unknown aggregate state**.
- Type of aggregate uncertainty matters for [▶ Surplus uncertainty](#)
 - how information affects trade volume.
 - what is optimal trading mechanism.
- Provision of incomplete information might harm welfare.
- Some implications find support in empirical literature.
[▶ Details](#)
- Extensions: entry and more general trading mechanisms.
[▶ Entry](#) [▶ Lotteries](#)

Appendix

Model predictions

Model predictions and evidence [◀ Go back](#)

- Bond markets
 - Increase in transparency → decrease in price dispersion: municipal bonds in US (Schultz, 2012).
 - More risky markets → more dispersed prices: OTC corporate bonds in US (Jankowitsch et al., 2011, Uslu and Velioglu, 2021).
 - Decrease in transparency → more market freezes possible; in contrast to Zou (2021).
 - States more different → more market freezes; as in Chiu and Koeppl, (2016).
- Labour markets
 - Increase in uncertainty → decrease in hiring; as in den Haan, Freund, Rendahl (2021).

US corporate bonds: dispersion and beta

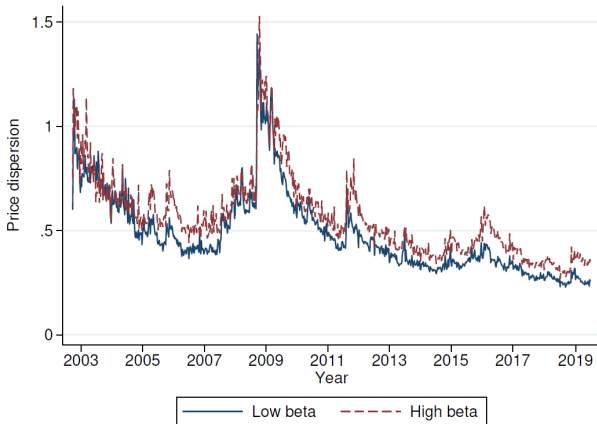


Figure: Uslu and Velioglu (2021)

US corporate bonds: dispersion and rating

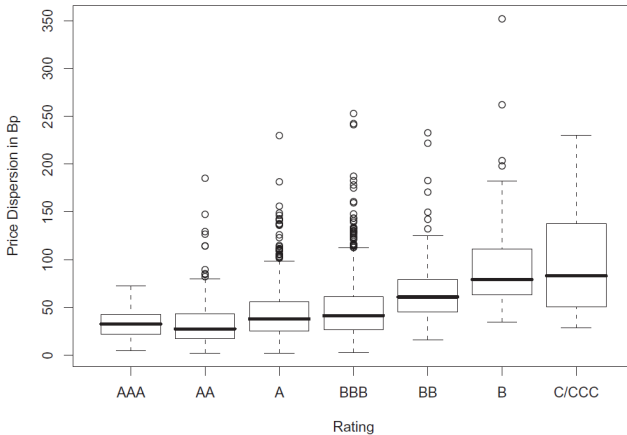


Figure: Jankowitsch et al (2011)

Equilibrium

Equilibrium conditions

- Focus on symmetric Nash equilibria

$$\left\{ \left\{ \kappa^{j,k}, x_i^{j,k}, p^{j,k} \right\}_{k=1}^{K^j}, U_i^j \right\}_j, \text{ for } i \in \{L, H\}, j \in \{B, G\}.$$

- Submarkets indexed by prices $p^{j,k}$.
- **Buyers** in state i for signal j **indifferent** between $p^{j,k}$:

$$U_i^{j,k} = (1 - \lambda(x_i^{j,k})) (v_i - p^{j,k}) \quad \text{for all } x_i^{j,k} > 0. \quad (3)$$

- Buyer-seller ratios consistent with measures of sellers and buyers: in state i , if $\sum_k^{K^j} x_i^{j,k} > 0$,

$$\sum_k^{K^j} \kappa^{j,k} x_i^{j,k} = \mathcal{B}_i. \quad (4)$$

Definition (Equilibrium)

We will say that a tuple $\left\{ \left\{ \kappa^{j,k}, x_i^{j,k}, p^{j,k} \right\}_{k=1}^{K^j}, U_i^j \right\}_j$ is an equilibrium for exogenous parameters $\Theta = (v_i, \mathcal{B}_i)$ with $i \in \{H, L\}$ and signal precision μ if for each $j \in \{G, B\}$:

- 1 given market utilities U_i^j , a tuple $\{x_i^{j,k}, p^{j,k}\}$ solves (1) and (2) for each k ,
- 2 market utilities satisfy (3) given $\{x_i^{j,k}, p^{j,k}\}_{k=1}^{K^j}$,
- 3 buyer-seller ratios and probability weights $\{x_i^{j,k}, \kappa^{j,k}\}_{k=1}^{K^j}$ are consistent with (4).

Pricing for both states

Lemma

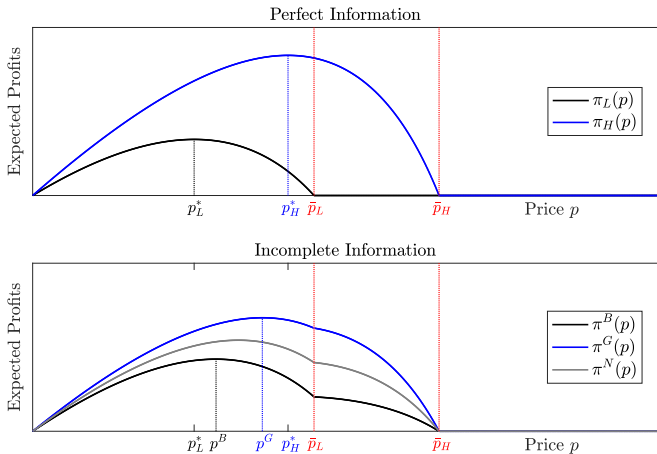
If the constrained profit maximisation on $[0, \min_i \bar{p}_i]$ has an interior solution p^j , then unique buyer-seller ratios $x_i^j > 0$ exist which together with p^j jointly solve:

$$\left(1 - \lambda(x_i^j)\right) (v_i - p^j) = U_i^j, \text{ for } i \in \{L, H\}, j \in \{B, G\}, \text{ and}$$

$$\mu \lambda(x_H^G) \left[\frac{\phi(x_H^G) v_H - p^G}{\phi(x_H^G) (v_H - p^G)} \right] + (1 - \mu) \lambda(x_L^G) \left[\frac{\phi(x_L^G) v_L - p^G}{\phi(x_L^G) (v_L - p^G)} \right] = 0,$$

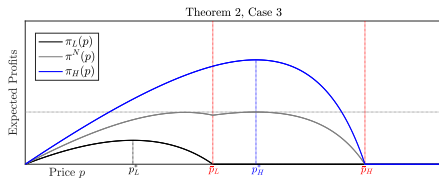
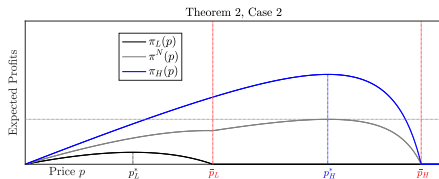
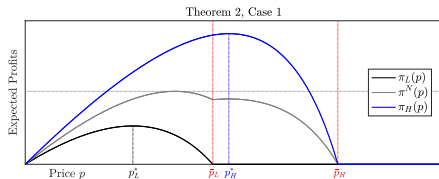
for $j = G$ and analogously for $j = B$. ▶ Simplifying assumptions

Small demand difference



Large demand difference

◀ Go back



Entry

Entry

Cost of setting up a trading post c :

- First best: entry depends on *demand state* and not signal.
- Second best: entry depends on signal, “pricing” as in perfect information, planner maximizes expected trades net of setup costs:

$$\mu\lambda(x_H^G)\phi(x_H^G)v + (1 - \mu)\lambda(x_L^G)\phi(x_L^G)v = c,$$

$$(1 - \mu)\lambda(x_H^B)\phi(x_H^B)v + \mu\lambda(x_L^B)\phi(x_L^B)v = c.$$

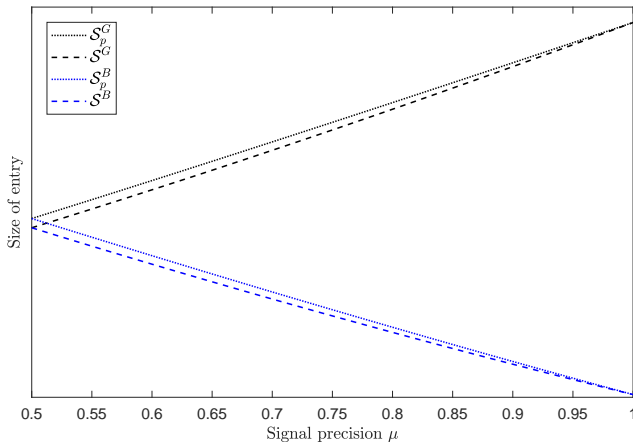
- Free entry conditions:

$$\left[\mu\lambda(x_H^G) + (1 - \mu)\lambda(x_L^G) \right] p^G = c,$$

$$\left[(1 - \mu)\lambda(x_H^B) + \mu\lambda(x_L^B) \right] p^B = c.$$

Information distorts entry

◀ Go back



Entry and information

Simplifying assumption

Assumption (Particular matching function)

Let $M(\mathcal{B}, \mathcal{S}) = \frac{\mathcal{B}\mathcal{S}}{\mathcal{B} + \mathcal{S}}$. Then $\lambda(x) = \phi(x) = \frac{x}{1+x}$, $\eta(x) = 1 - \lambda(x)$.

Assumption (Particular matching function)

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- Under tightness risk: for $i \in \{L, H\}$ and $j = G$

$$\left(1 - \lambda\left(x_i^G\right)\right) \left(v - p^G\right) = U_i^G, \text{ and}$$

$$\mu\lambda\left(x_H^G\right) \left[\frac{\lambda\left(x_H^G\right)v - p^G}{\lambda\left(x_H^G\right)\left(v - p^G\right)}\right] + (1 - \mu)\lambda\left(x_L^G\right) \left[\frac{\lambda\left(x_L^G\right)v - p^G}{\lambda\left(x_L^G\right)\left(v - p^G\right)}\right] = 0.$$

◀ Go back

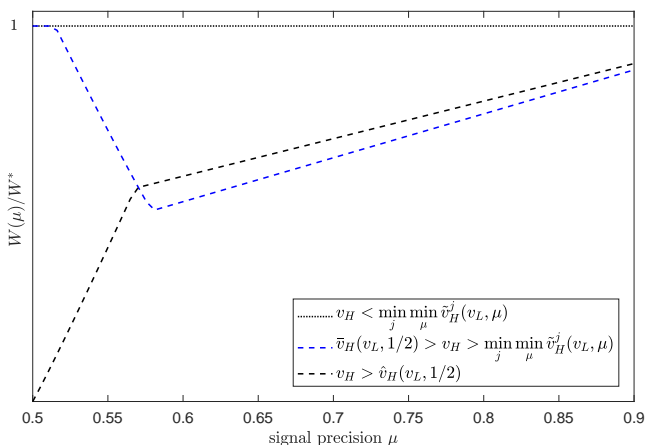
Surplus risk

Surplus risk

- If $v_H - v_L$ is small, only single price equilibria exist.
- If $v_H - v_L$ is large, equilibrium under no information features $p^N = \phi(\mathcal{B})v_H > v_L$ and no trade in L -state:
 - interval $(\underline{\mu}, \bar{\mu})$ exists where sellers mix and expected value of trades is greater than under no information.
 - qualitatively similar to tightness risk otherwise.
- Numerically find thresholds for each type of equilibrium.

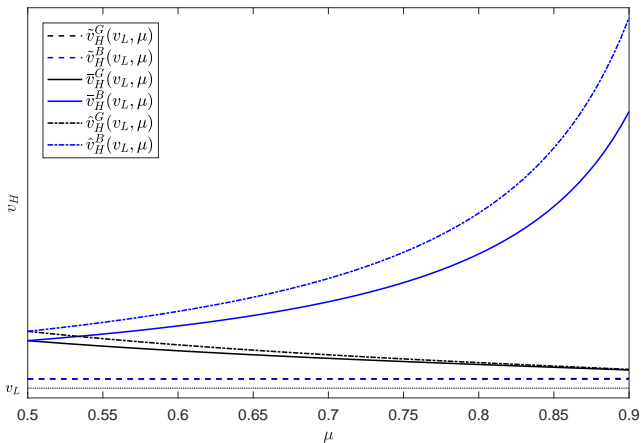
Information and welfare: surplus risk

◀ Go back



Normalised welfare for three different v_H (fixed v_L and \mathcal{B}).

Existence thresholds: surplus risk



Lotteries

- Suppose sellers can post lotteries (θ_i^j, p_i^j) :
 - lottery price p_i^j , and
 - probability of getting the good θ_i^j .
- Lotteries screen states if

$$\theta_i^j v_i - p_i^j \geq \theta_{-i}^j v_i - p_{-i}^j.$$

- We show under
 - tightness risk, price posting is optimal: $\theta_i^j = 1$ and $p_i^j = p^j$.
 - surplus risk, lotteries weakly dominate price posting: zero trade in low state no longer equilibrium outcome.